

# IP Forwarding

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Covers the principles of end-to-end datagram delivery in IP networks.

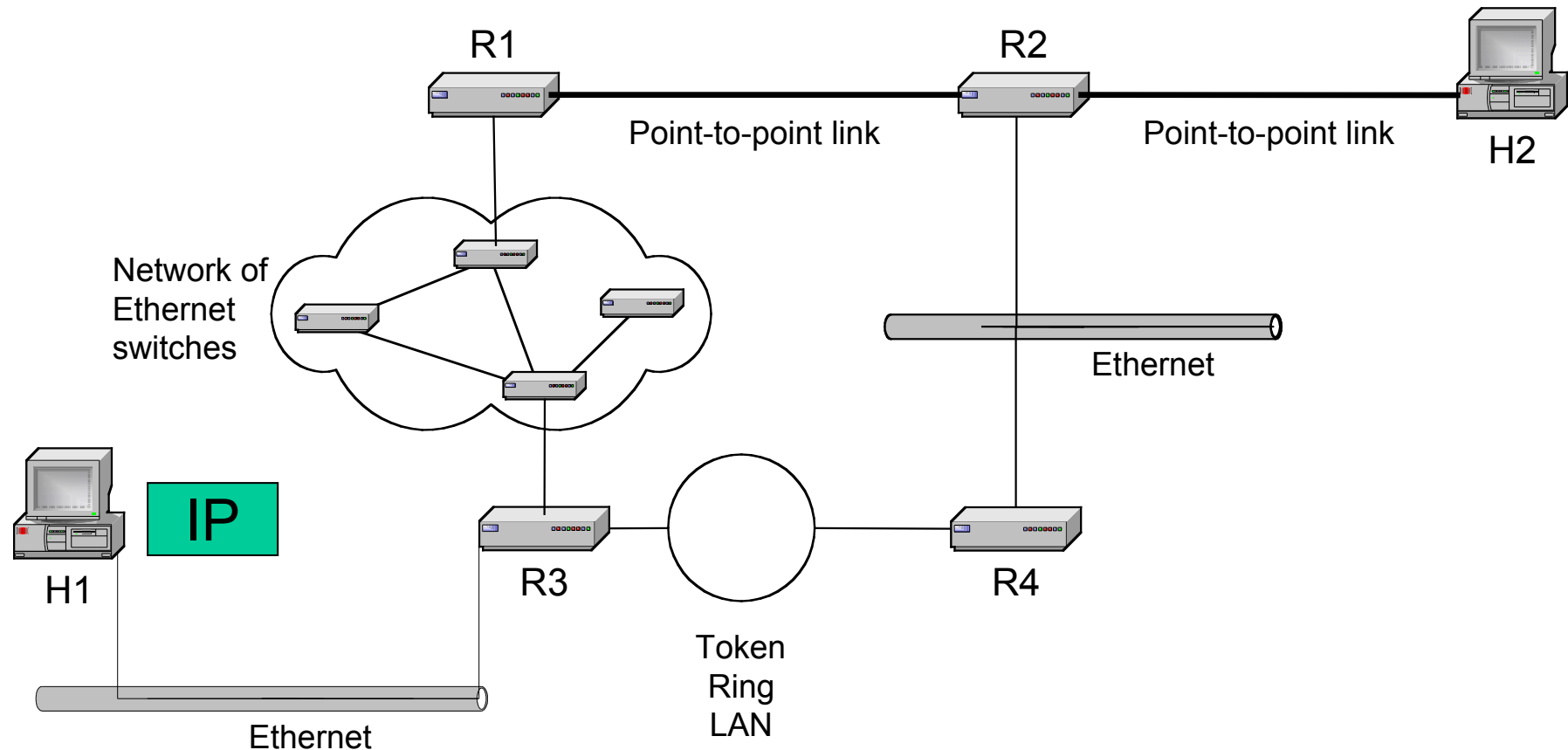
# Orientation

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- Internet is a collection of networks
- IP provides an end-to-end delivery service for IP datagrams between hosts
- The delivery service is realized with the help of IP routers

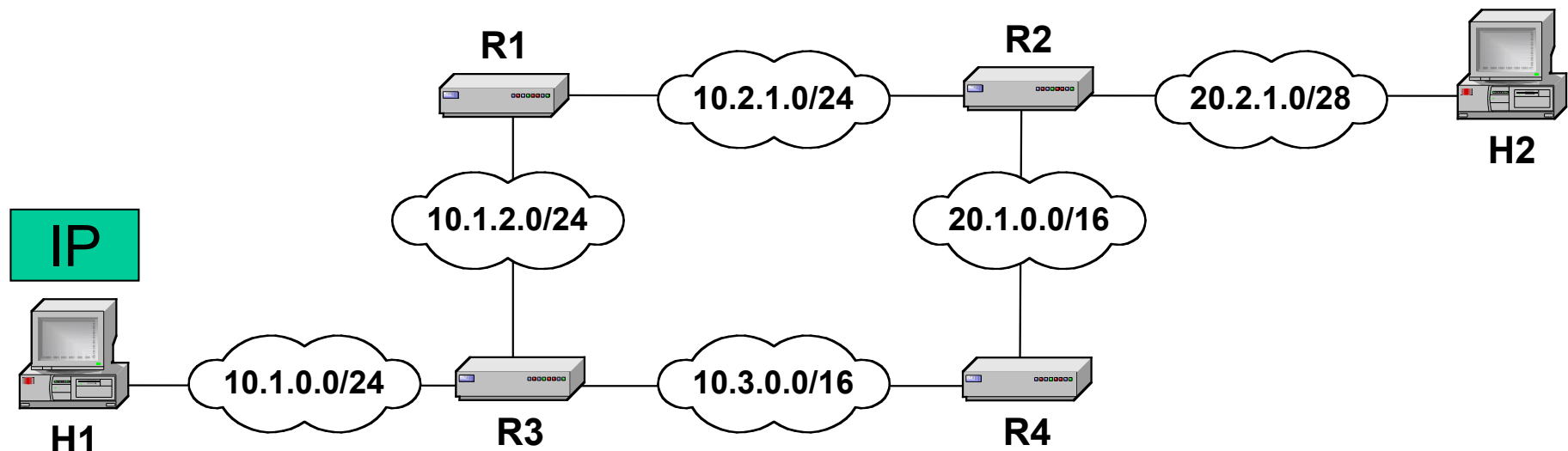
# Delivery of an IP datagram

- View at the data link layer layer:
  - Internetwork is a collection of LANs or point-to-point links or switched networks that are connected by routers



# Delivery of an IP datagram

- View at the IP layer:
  - An IP network is a logical entity with a network number
  - We represent an IP network as a “cloud”
  - The IP delivery service takes the view of clouds, and ignores the data link layer view



# Tenets of end-to-end delivery of datagrams

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The following conditions must hold so that an IP datagram can be successfully delivered

1. The network prefix of an IP destination address must correspond to a unique data link layer network (=LAN or point-to-point link or switched network).  
(The reverse need not be true!)
2. Routers and hosts that have a common network prefix must be able to exchange IP datagrams using a data link protocol (e.g., Ethernet, PPP)
3. Every data link layer network must be connected to at least one other data link layer network via a router.


# Routing tables

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- Each router and each host keeps a **routing table** which tells the router how to process an outgoing packet
- Main columns:
  1. **Destination address:** where is the IP datagram going to?
  2. **Next hop or interface:** how to send the IP datagram?
- Routing tables are set so that a datagram gets closer to the its destination every hop

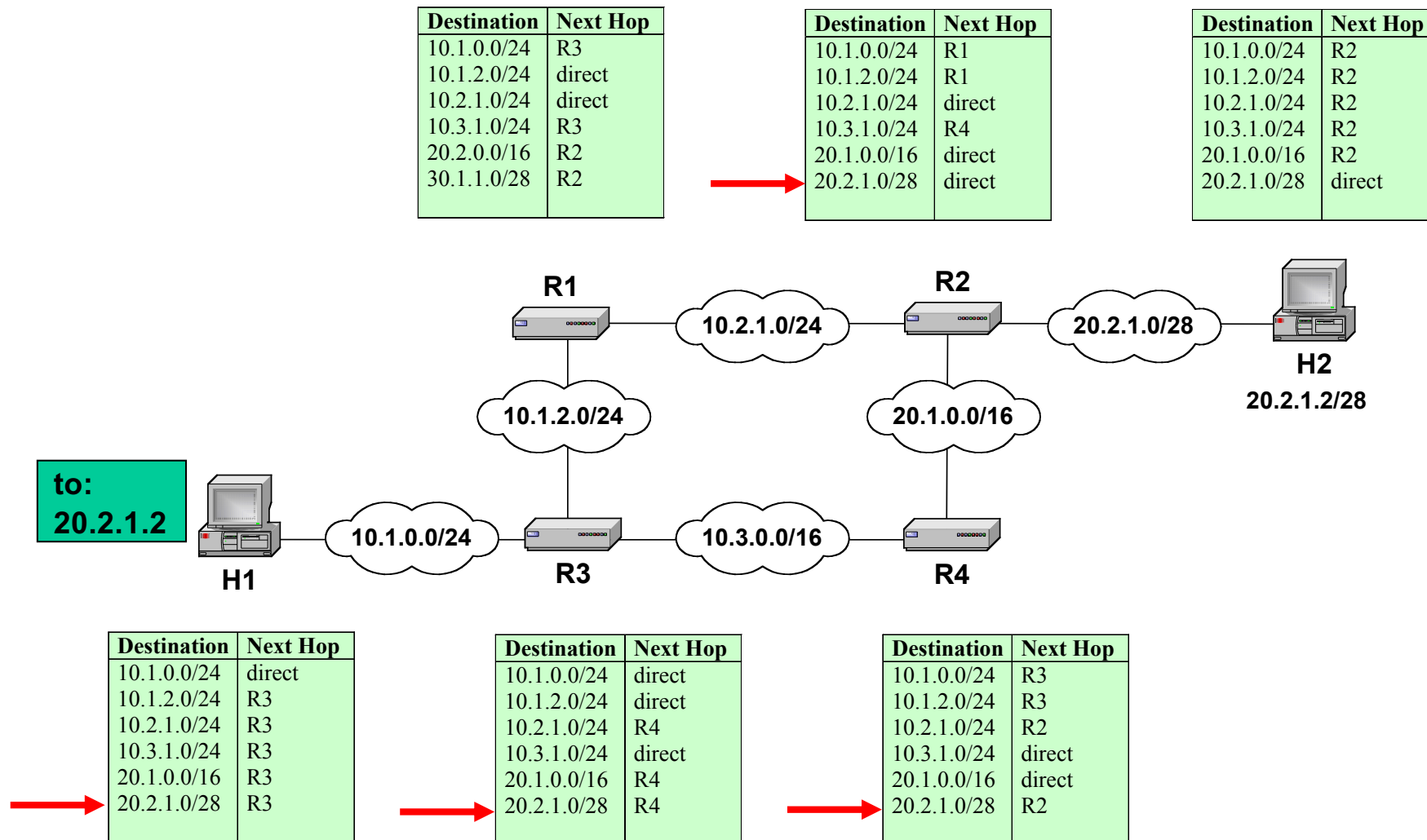
Routing table of a host or router

IP datagrams can be directly delivered (“direct”) or are sent to a router (“R4”)



Destination	Next Hop
10.1.0.0/24	direct
10.1.2.0/24	direct
10.2.1.0/24	R4
10.3.1.0/24	direct
20.1.0.0/16	R4
20.2.1.0/28	R4

# Delivery with routing tables



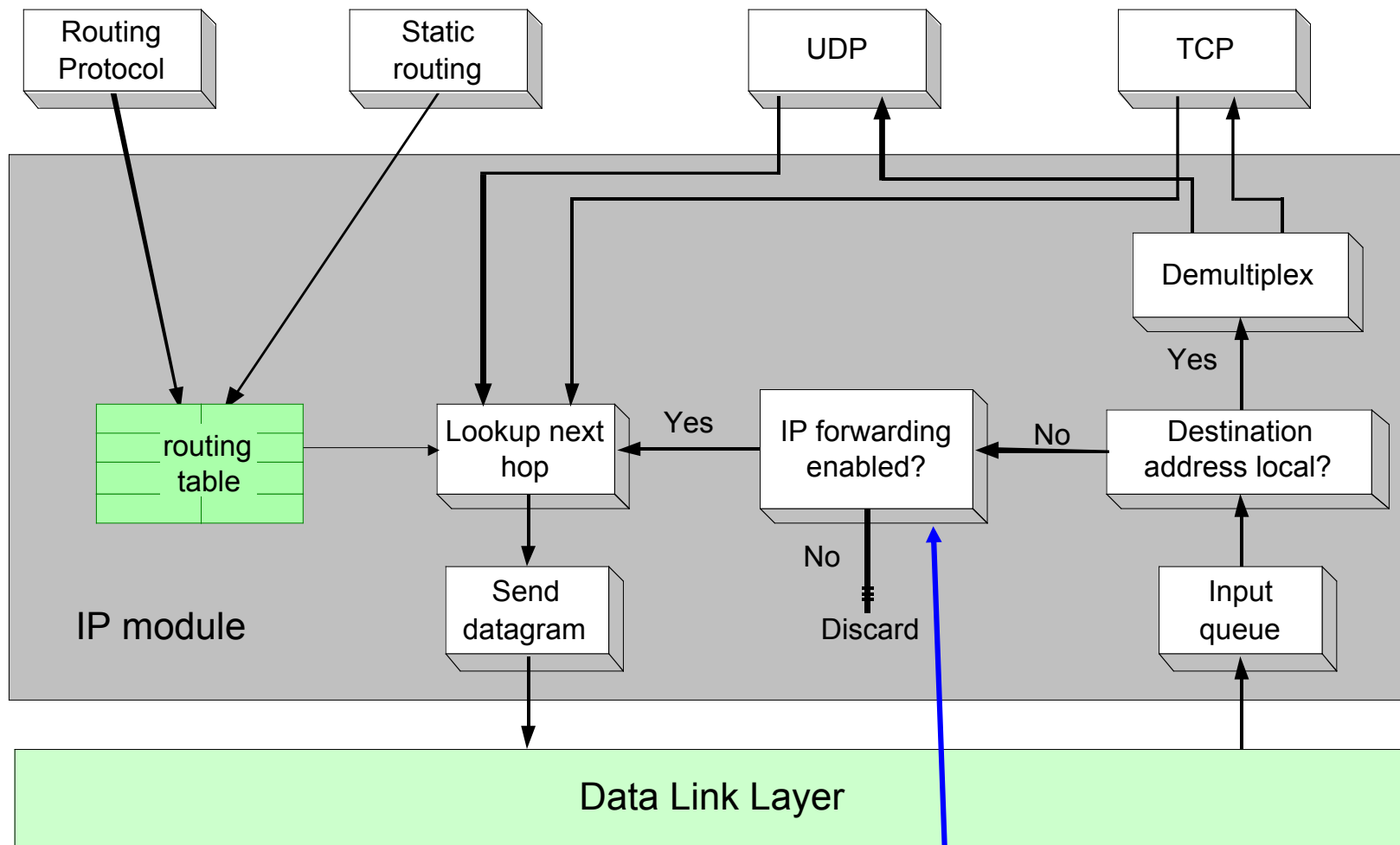
# Delivery of IP datagrams

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- There are two distinct processes to delivering IP datagrams:
  1. **Forwarding:** How to pass a packet from an input interface to the output interface?
  2. **Routing:** How to find and setup the routing tables?
- Forwarding must be done as fast as possible:
  - on routers, is often done with support of hardware
  - on PCs, is done in kernel of the operating system
- Routing is less time-critical
  - On a PC, routing is done as a background process



# Processing of an IP datagram in IP



IP router: IP forwarding enabled  
Host: IP forwarding disabled 9

# Processing of an IP datagram in IP

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- Processing of IP datagrams is very similar on an IP router and on a host
- **Main difference:**  
“IP forwarding” is enabled on router and disabled on host
- **IP forwarding enabled**  
→ if a datagram is received, but it is not for the local system, the datagram will be sent to a different system
- **IP forwarding disabled**  
→ if a datagram is received, but it is not for the local system, the datagram will be discarded

# Processing of an IP datagram at a router

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**Receive an  
IP datagram**



1. IP header validation
2. Process options in IP header
3. Parsing the destination IP address
4. Routing table lookup
5. Decrement TTL
6. Perform fragmentation (if necessary)
7. Calculate checksum
8. Transmit to next hop
9. Send ICMP packet (if necessary)

# Routing table lookup

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- When a router or host needs to transmit an IP datagram, it performs a routing table lookup
- **Routing table lookup:** Use the IP destination address as a key to search the routing table.
- Result of the lookup is the IP address of a next hop router, or the name of a network interface

Destination address	Next hop
network prefix <i>or</i> host IP address <i>or</i> loopback address <i>or</i> default route	IP address of next hop router  <i>or</i>  Name of a network interface

# Type of routing table entries

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- **Network route**
  - Destination address is a network address (e.g., 10.0.2.0/24)
  - Most entries are network routes
- **Host route**
  - Destination address is an interface address (e.g., 10.0.1.2/32)
  - Used to specify a separate route for certain hosts
- **Default route**
  - Used when no network or host route matches
  - The router that is listed as the next hop of the default route is the **default gateway (for Cisco: “gateway of last resort)**
- **Loopback address**
  - Loopback address: 127.0.0.1
  - The next hop lists the loopback (lo0) interface as outgoing interface

# Longest Prefix Match

- **Longest Prefix Match:** Search for the routing table entry that has the longest match with the prefix of the destination IP address

1. Search for a match on all 32 bits
2. Search for a match for 31 bits
- .....
32. Search for a mach on 0 bits

Host route, loopback entry  
→ 32-bit prefix match

Default route is represented as 0.0.0.0/0  
→ 0-bit prefix match

128.143.71.21



Destination address	Next hop
10.0.0.0/8	R1
128.143.0.0/16	R2
128.143.64.0/20	R3
128.143.192.0/20	R3
128.143.71.0/24	R4
128.143.71.55/32	R3
default	R5



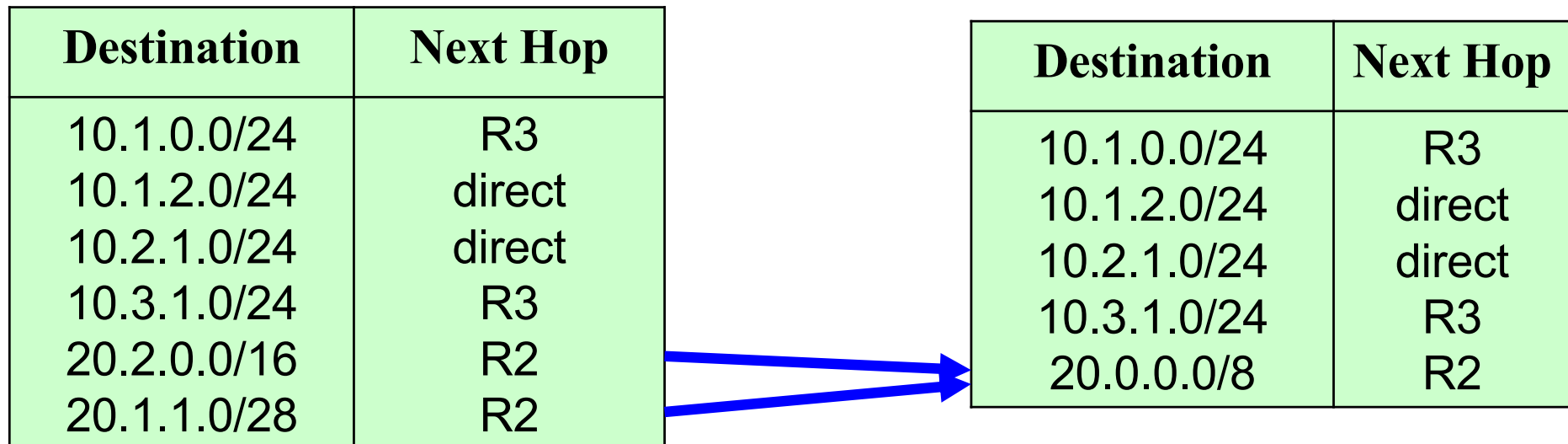
**The longest prefix match for 128.143.71.21 is for 24 bits with entry 128.143.71.0/24**

**Datagram will be sent to R4**

# Route Aggregation

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- Longest prefix match algorithm permits to aggregate prefixes with **identical next hop address** to a single entry
- This contributes significantly to reducing the size of routing tables of Internet routers



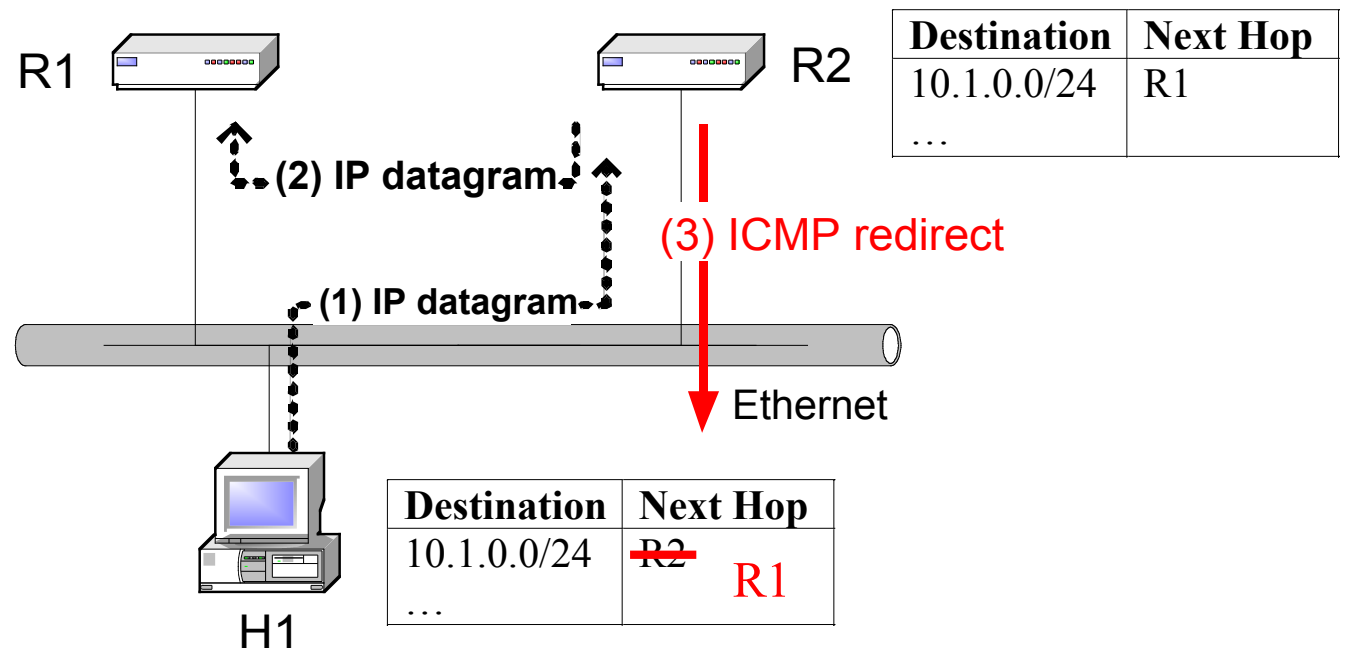
The diagram illustrates the process of route aggregation. On the left, a routing table with 6 entries is shown. Two blue arrows point from the last two entries of this table to the last entry of a second routing table on the right. This second table shows the result of aggregating the last two entries of the first table into a single entry with a more specific destination prefix.

Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	direct
10.2.1.0/24	direct
10.3.1.0/24	R3
20.2.0.0/16	R2
20.1.1.0/28	R2

Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	direct
10.2.1.0/24	direct
10.3.1.0/24	R3
20.0.0.0/8	R2

# Routing table manipulations with ICMP

- When a router detects that an IP datagram should have gone to a different router, the router (here R2)
  - forwards the IP datagram to the correct router
  - sends an ICMP redirect message to the host
- Host uses ICMP message to update its routing table





# ICMP Router Solicitation

# ICMP Router Advertisement

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## Router Discovery Protocol

1. After bootstrapping, a Host broadcasts or multicast (address 224.0.0.2) an **ICMP Router Solicitation** Messages.
2. In response, Routers send an **ICMP router advertisement** message
3. Also, routers periodically broadcast **ICMP router advertisement**
4. The host chooses the **first message** it receives and adds that router to its routing table.

