

IP Addressing

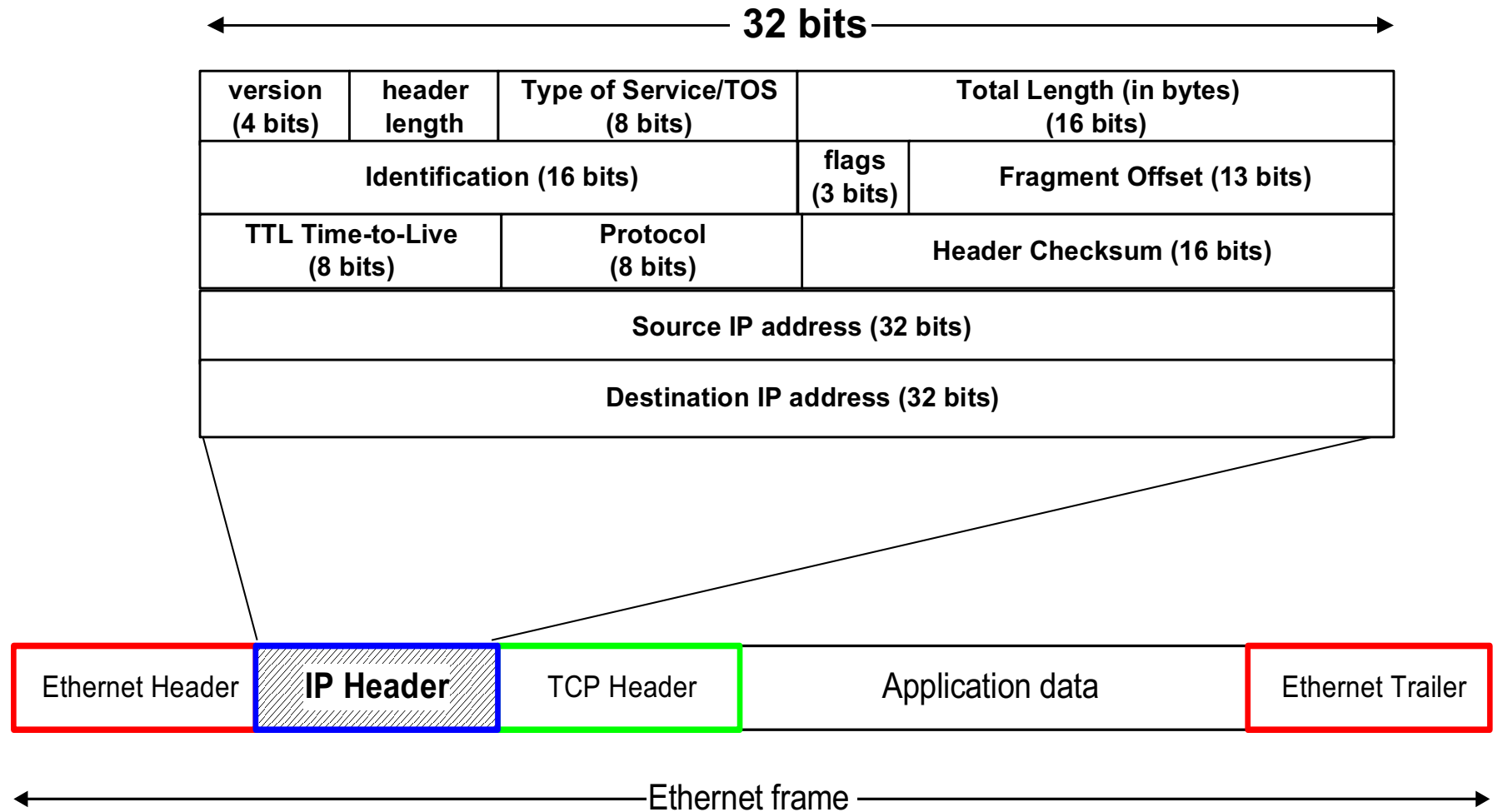
Introductory material.

An entire module devoted to IP addresses.

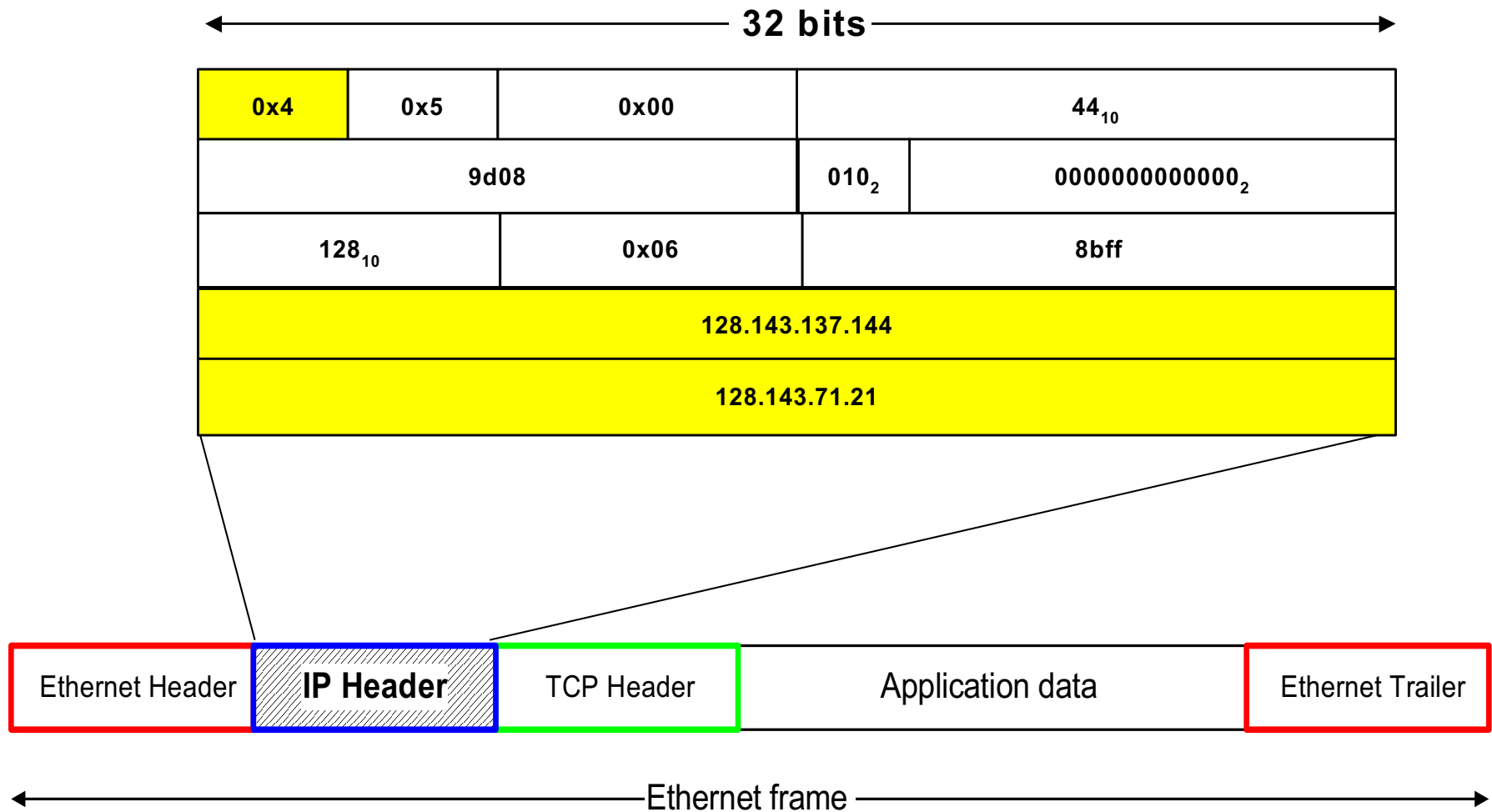
IP Addresses

- Structure of an IP address
- Classful IP addresses
- Limitations and problems with classful IP addresses
- Subnetting
- CIDR
- IP Version 6 addresses

IP Addresses



IP Addresses



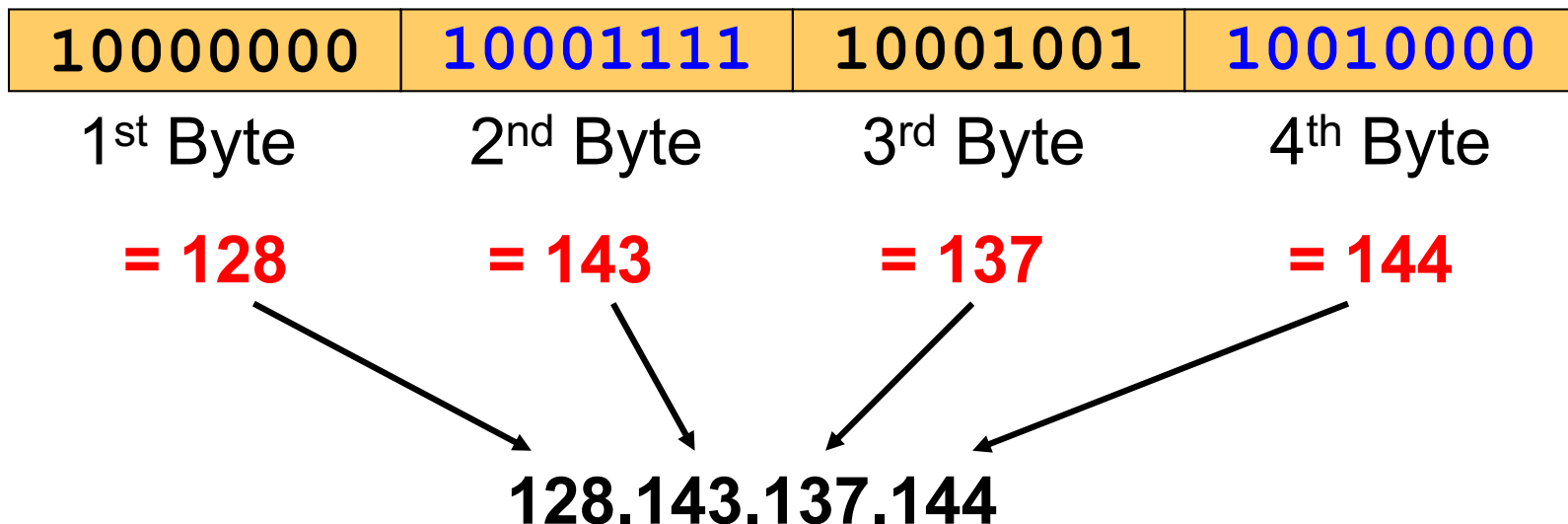
What is an IP Address?

- An IP address is a unique global address for a network interface
- An IP address:
 - is a **32 bit long** identifier
 - encodes a network number (**network prefix**) and a **host number**

Dotted Decimal Notation

- IP addresses are written in a so-called *dotted decimal notation*
- Each byte is identified by a decimal number in the range [0..255]:

- **Example:**



Network prefix and Host number

- The network prefix identifies a network and the host number identifies a specific host (actually, interface on the network).

network prefix

host number

- **How do we know how long the network prefix is?**
 - The network prefix is implicitly defined (see **class-based addressing**)
 - The network prefix is indicated by a **netmask.**

Example

- **Example:** ellington.cs.virginia.edu

128.143

137.144

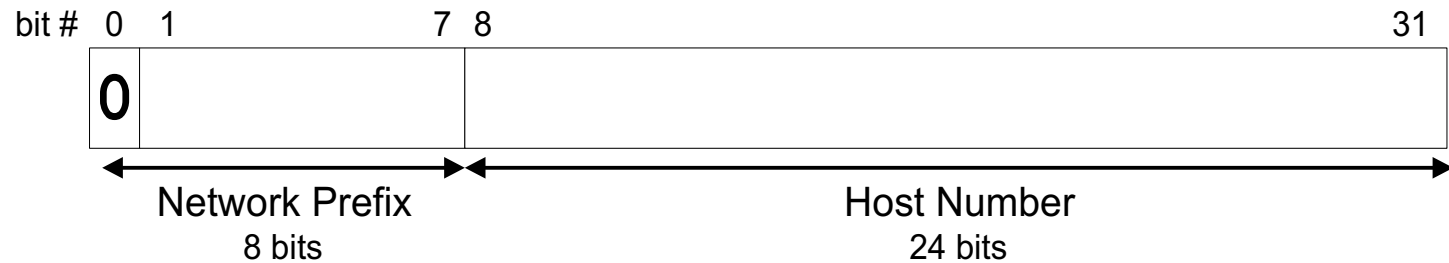
- Network id is: 128.143.0.0
- Host number is: 137.144
- Network mask is: 255.255.0.0 or `ffff0000`
- Prefix notation: **128.143.137.144/16**
 - » Network prefix is 16 bits long

The old way: Classful IP Addresses

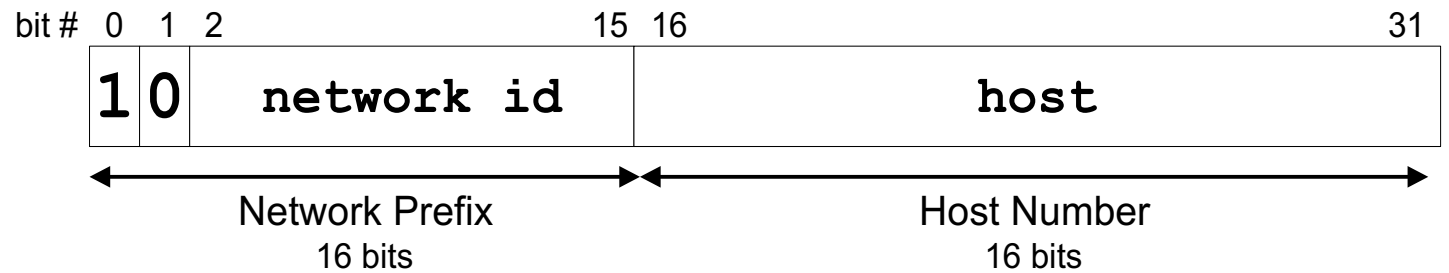
- When Internet addresses were standardized (early 1980s), the Internet address space was divided up into classes:
 - **Class A:** Network prefix is 8 bits long
 - **Class B:** Network prefix is 16 bits long
 - **Class C:** Network prefix is 24 bits long
- Each IP address contained a key which identifies the class:
 - **Class A:** IP address starts with “0”
 - **Class B:** IP address starts with “10”
 - **Class C:** IP address starts with “110”

The old way: Internet Address Classes

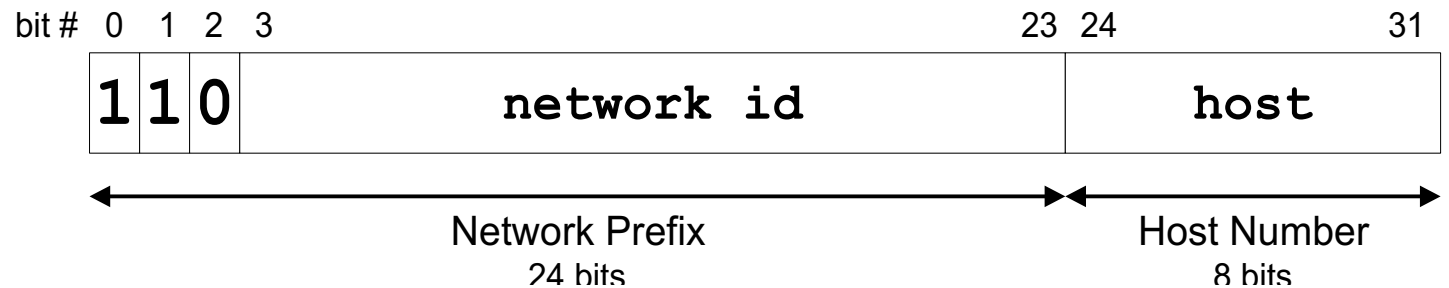
Class A



Class B



Class C



Problems with Classful IP Addresses

- The original classful address scheme had a number of problems

Problem 1. Too few network addresses for large networks

- Class A and Class B addresses are gone

Problem 2. Two-layer hierarchy is not appropriate for large networks with Class A and Class B addresses.

- **Fix #1: Subnetting**

Problems with Classful IP Addresses

Problem 3. Inflexible. Assume a company requires 2,000 addresses

- Class A and B addresses are overkill
- Class C address is insufficient (requires 8 Class C addresses)
- **Fix #2: Classless Interdomain Routing (CIDR)**

Problems with Classful IP Addresses

Problem 4: Exploding Routing Tables: Routing on the backbone Internet needs to have an entry for each network address. In 1993, the size of the routing tables started to outgrow the capacity of routers.

– **Fix #2: Classless Interdomain Routing (CIDR)**

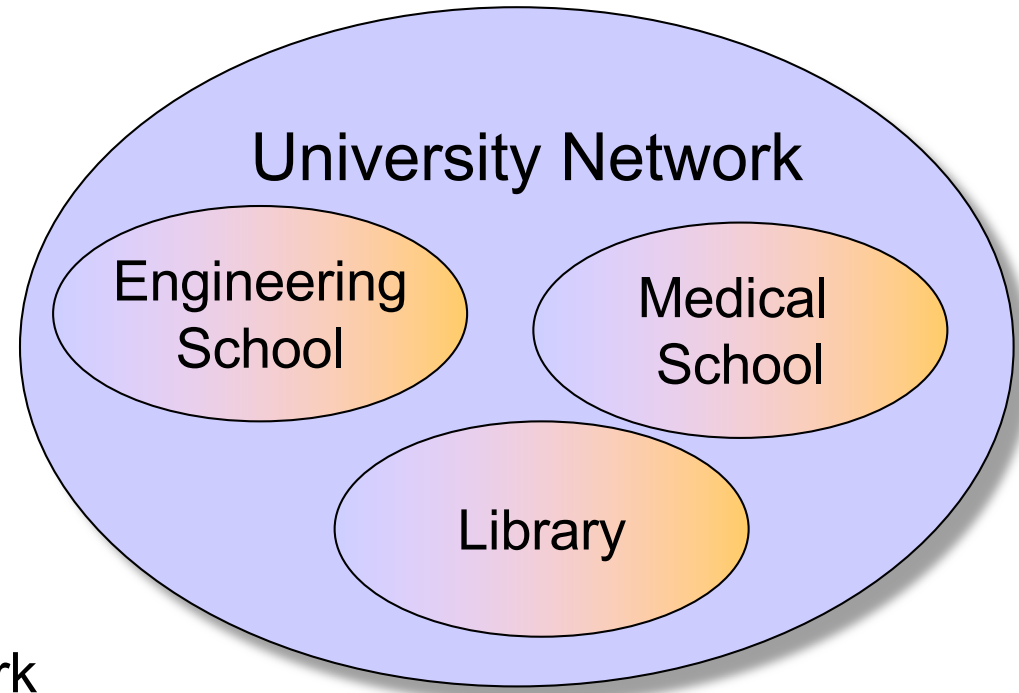
Problems with Classful IP Addresses

Problem 5. The Internet is going to outgrow the 32-bit addresses

– **Fix #3: IP Version 6**

Subnetting

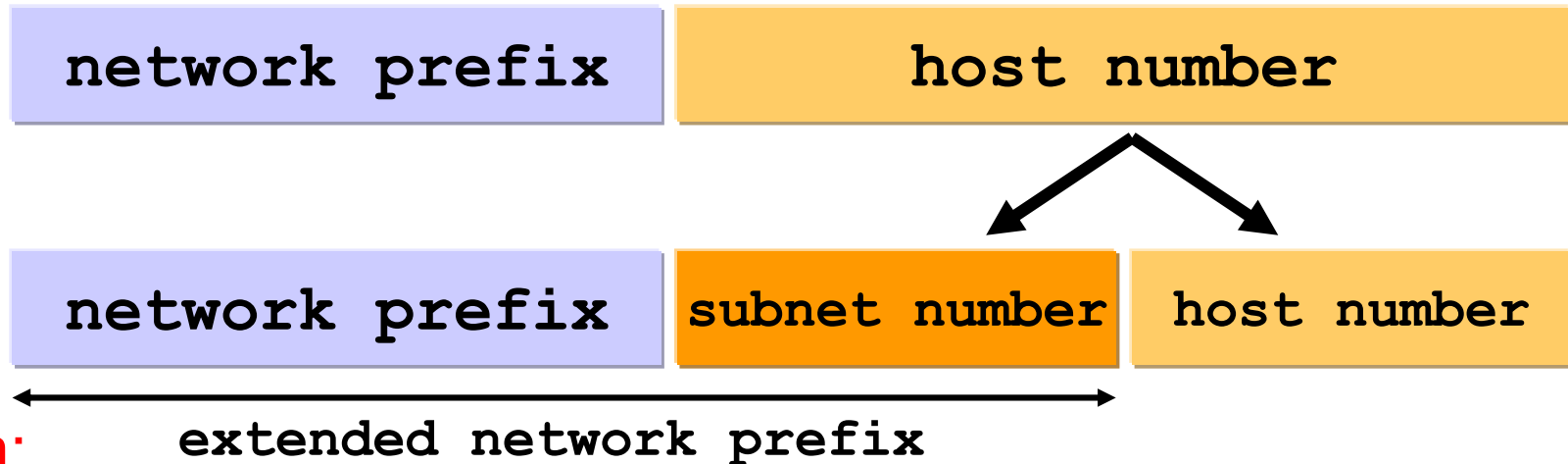
- **Problem:** Organizations have multiple networks which are independently managed
 - **Solution 1:** Allocate one or more Class C address for each network
 - Difficult to manage
 - From the outside of the organization, each network must be addressable.
 - **Solution 2:** Add another level of hierarchy to the IP addressing structure



→ Subnetting

Basic Idea of Subnetting

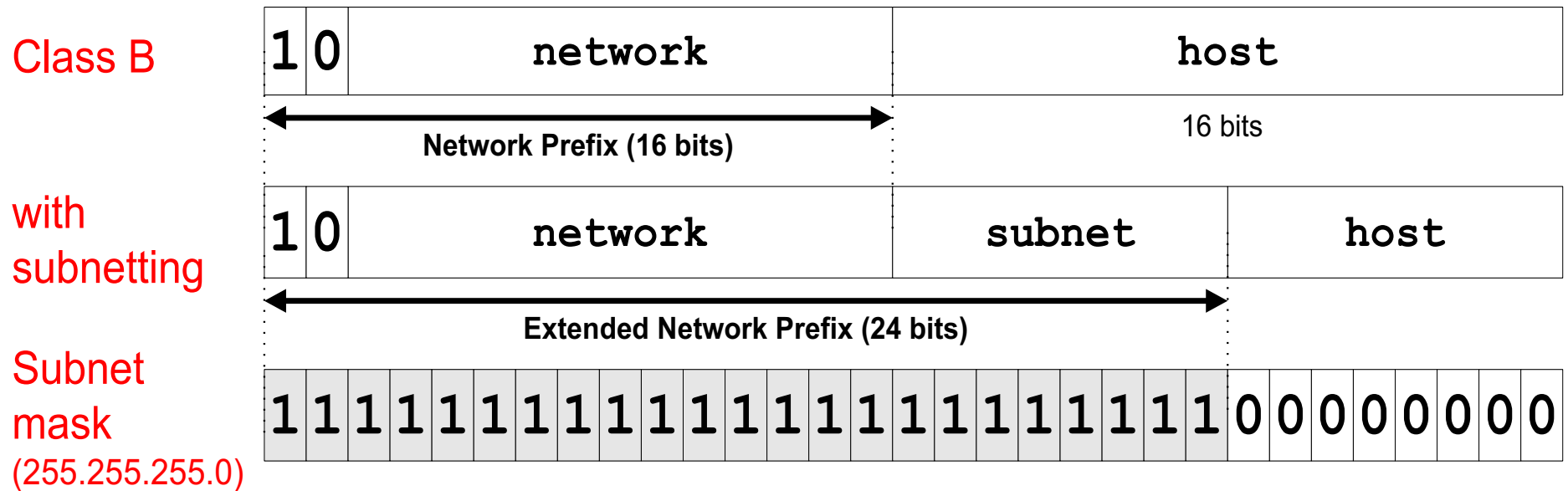
- Split the host number portion of an IP address into a **subnet number** and a (smaller) **host number**.
- Result is a 3-layer hierarchy



- **Then:**
 - Subnets can be freely assigned within the organization
 - Internally, subnets are treated as separate networks
 - Subnet structure is not visible outside the organization

Subnet Masks

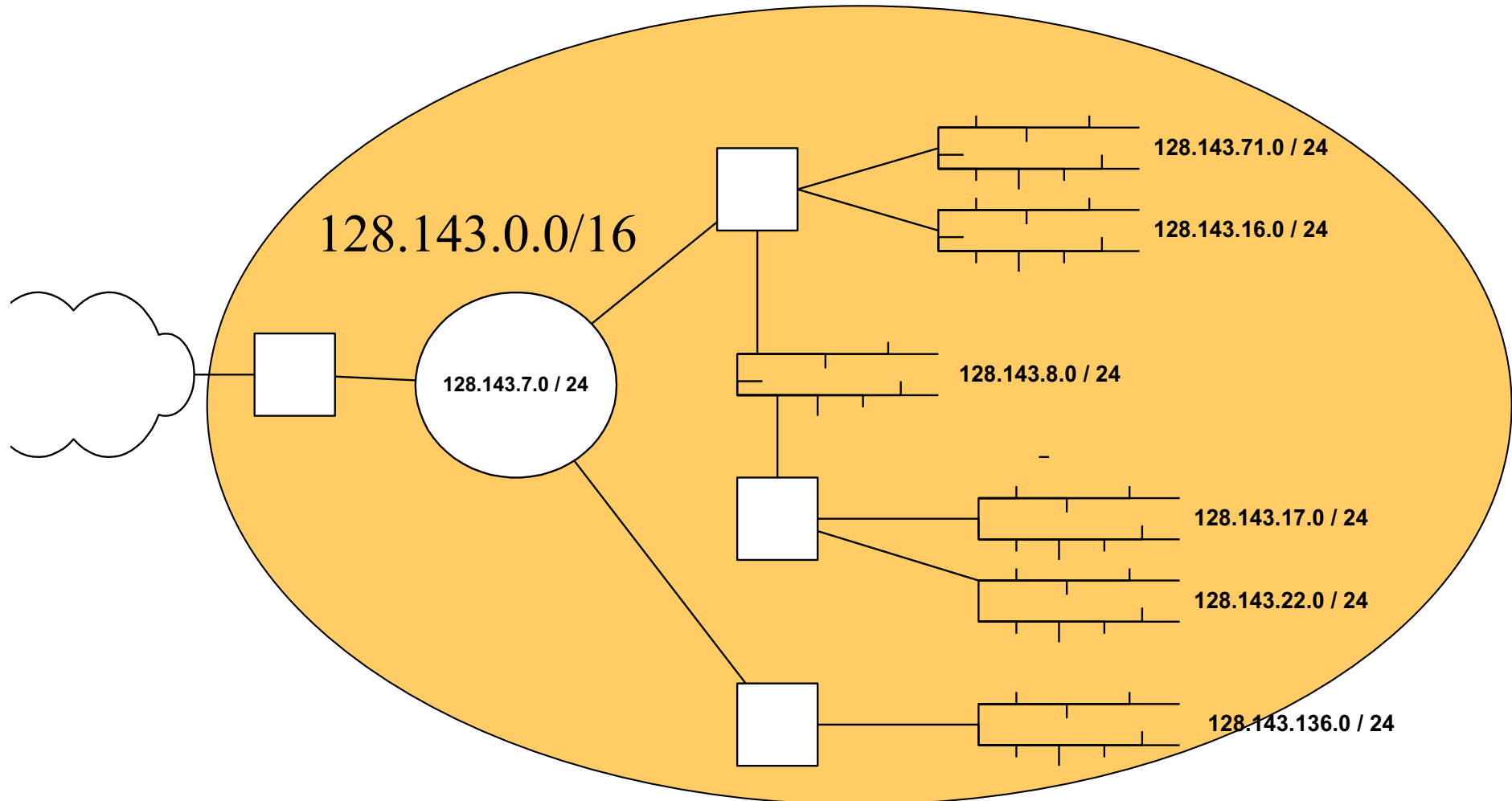
- Routers and hosts use an **extended network prefix (subnet mask)** to identify the start of the host numbers



- * There are different ways of subnetting. Commonly used netmasks for university networks with /16 prefix (Class B) are 255.255.255.0 and 255.255.0.0

Typical Addressing Plan for an Organization that uses subnetting

- Each layer-2 network (Ethernet segment, FDDI segment) is allocated a subnet address.



Advantages of Subnetting

- With subnetting, IP addresses use a 3-layer hierarchy:
 - » Network
 - » Subnet
 - » Host
- Improves efficiency of IP addresses by not consuming an entire Class B or Class C address for each physical network/
- Reduces router complexity. Since external routers do not know about subnetting, the complexity of routing tables at external routers is reduced.
- Note: Length of the subnet mask need not be identical at all subnetworks.

CIDR - Classless Interdomain Routing

- IP backbone routers have one routing table entry for each network address:
 - With subnetting, a backbone router only needs to know one entry for each Class A, B, or C networks
 - This is acceptable for Class A and Class B networks
 - $2^7 = 128$ Class A networks
 - $2^{14} = 16,384$ Class B networks
 - But this is not acceptable for Class C networks
 - $2^{21} = 2,097,152$ Class C networks
- In 1993, the size of the routing tables started to outgrow the capacity of routers
- Consequence: The Class-based assignment of IP addresses had to be abandoned

CIDR - Classless Interdomain Routing

- **Goals:**
 - Restructure IP address assignments to increase efficiency
 - Hierarchical routing aggregation to minimize route table entries
- **CIDR (Classless Interdomain routing)** abandons the notion of classes:
Key Concept: The length of the network id (prefix) in the IP addresses is kept arbitrary
- **Consequence:** Routers advertise the IP address and the length of the prefix

CIDR Example

- CIDR notation of a network address:
192.0.2.0/18
 - "18" says that the first 18 bits are the network part of the address (and 14 bits are available for specific host addresses)
- The network part is called the **prefix**
- Assume that a site requires a network address with 1000 addresses
- With CIDR, the network is assigned a continuous block of 1024 addresses with a 22-bit long prefix

CIDR: Prefix Size vs. Network Size

CIDR Block Prefix	# of Host Addresses
/27	32 hosts
/26	64 hosts
/25	128 hosts
/24	256 hosts
/23	512 hosts
/22	1,024 hosts
/21	2,048 hosts
/20	4,096 hosts
/19	8,192 hosts
/18	16,384 hosts
/17	32,768 hosts
/16	65,536 hosts
/15	131,072 hosts
/14	262,144 hosts
/13	524,288 hosts

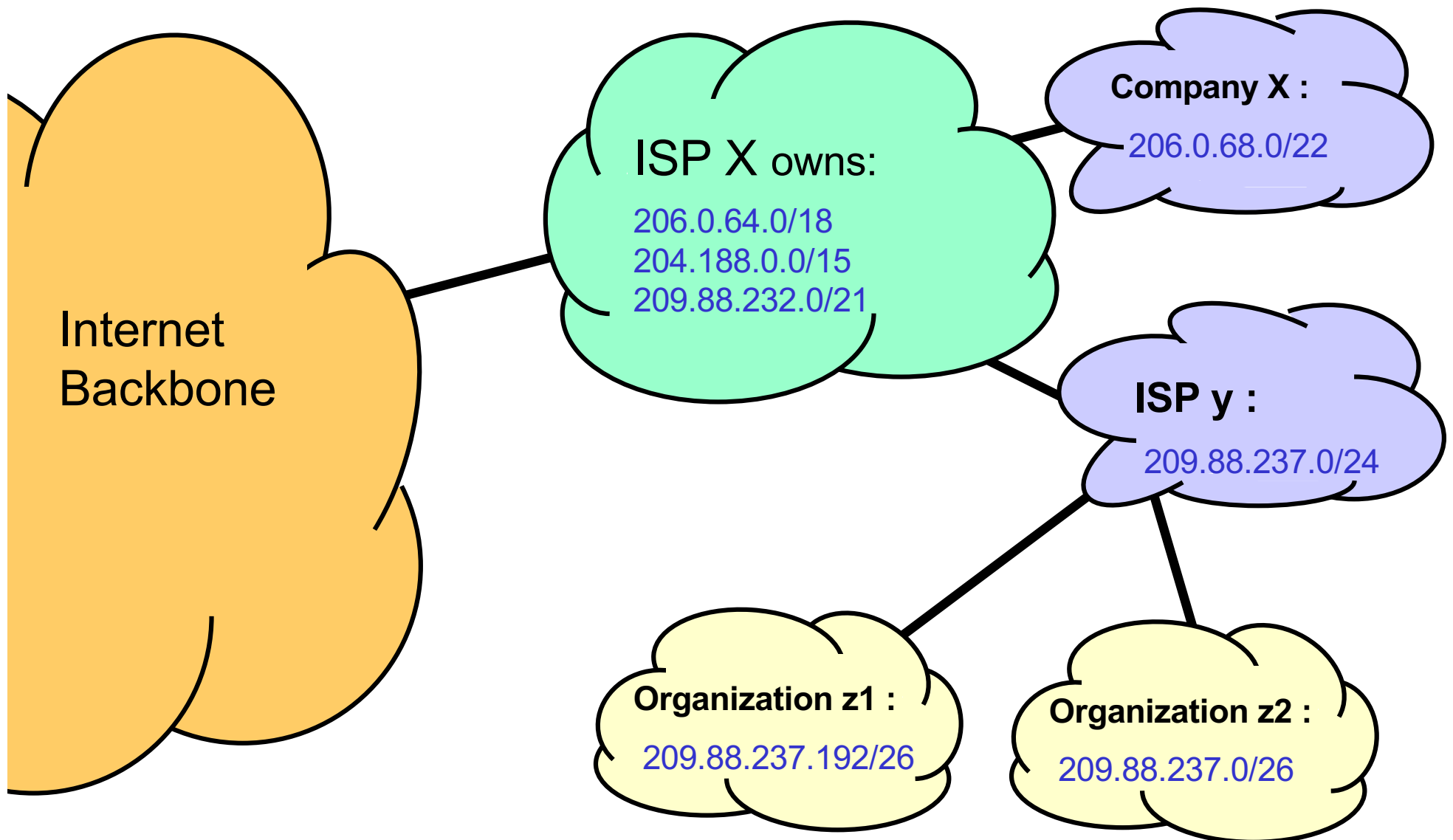
CIDR and Address assignments

- Backbone ISPs obtain large block of IP addresses space and then reallocate portions of their address blocks to their customers.

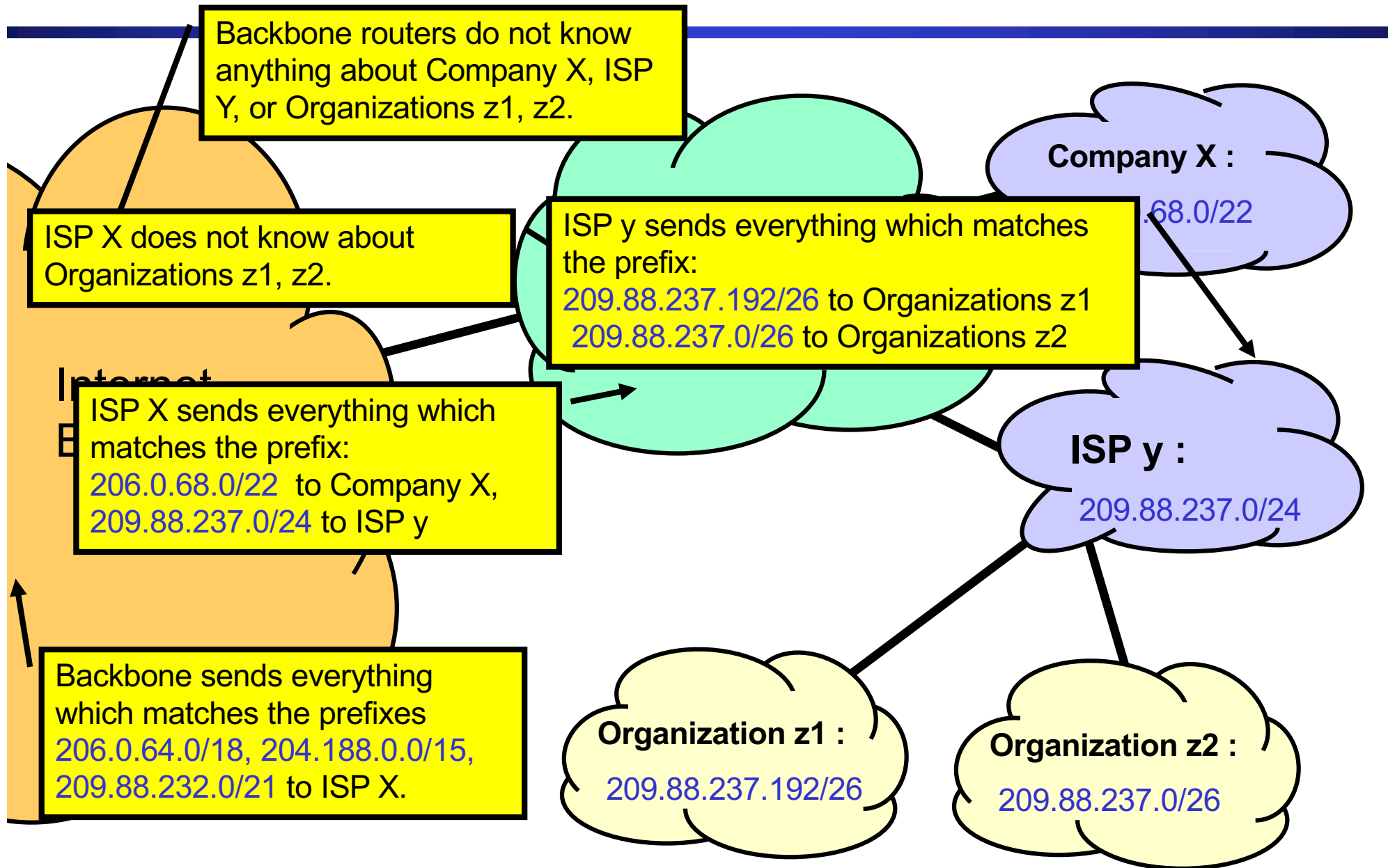
Example:

- Assume that an ISP owns the address block [206.0.64.0/18](#), which represents 16,384 (2^{14}) IP addresses
- Suppose a client requires 800 host addresses
- With classful addresses: need to assign a class B address (and waste ~64,700 addresses) or four individual Class Cs (and introducing 4 new routes into the global Internet routing tables)
- With CIDR: Assign a /22 block, e.g., [206.0.68.0/22](#), and allocated a block of 1,024 (2^{10}) IP addresses.

CIDR and Routing Information



CIDR and Routing Information



Example

You can find about ownership of IP addresses in North America via <http://www.arin.net/whois/>

- The IP Address: **207.2.88.170**

207

2

88

170

11001111	00000010	01011000	10101010
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Belongs to:

City of Charlottesville, VA: **207.2.88.0 - 207.2.92.255**

11001111	00000010	01011000	00000000
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Belongs to:

Cable & Wireless USA **207.0.0.0 - 207.3.255.255**

11001111	00000000	00000000	00000000
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CIDR and Routing

- **Aggregation of routing table entries:**
 - 128.143.0.0/16 and 128.144.0.0/16 are represented as 128.142.0.0/15
- **Longest prefix match:** Routing table lookup finds the routing entry that matches the the longest prefix

What is the outgoing interface for 128.143.137.0/24 ?

Prefix	Interface
128.0.0.0/4	interface #5
128.128.0.0/9	interface #2
128.143.128.0/17	interface #1

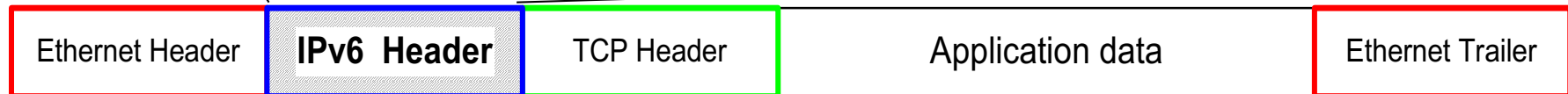
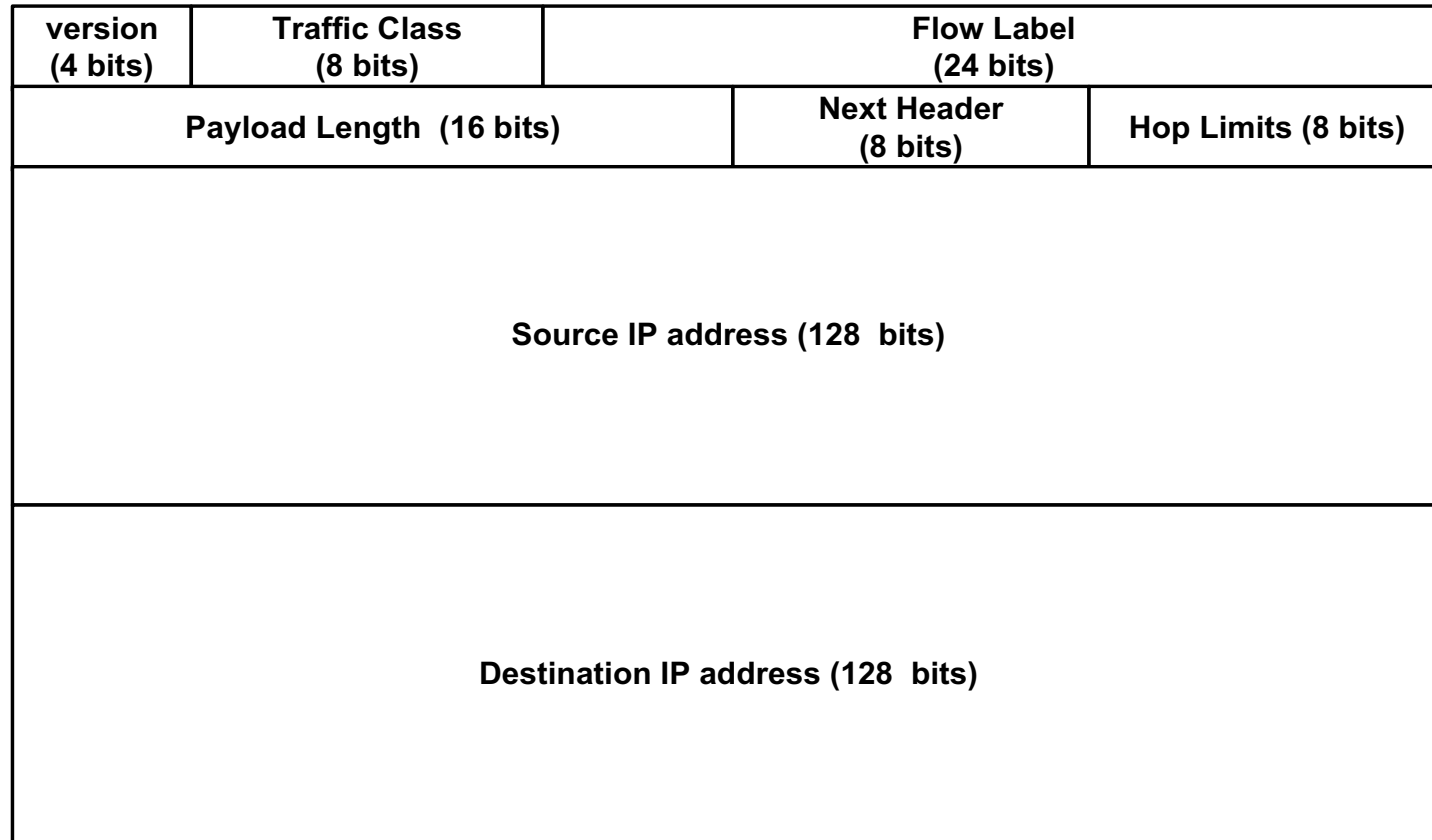
Routing table

IPv6 - IP Version 6

- **IP Version 6**
 - Is the successor to the currently used IPv4
 - Specification completed in 1994
 - Makes improvements to IPv4 (no revolutionary changes)
- One (not the only !) feature of IPv6 is a significant increase in of the IP address to **128 bits (16 bytes)**
 - IPv6 will solve – for the foreseeable future – the problems with IP addressing

IPv6 Header

32 bits



Ethernet frame

IPv6 vs. IPv4: Address Comparison

- **IPv4** has a maximum of
 $2^{32} \approx 4$ billion addresses
- **IPv6** has a maximum of
 $2^{128} = (2^{32})^4 \approx 4$ billion x 4 billion x 4 billion x 4 billion
addresses

Notation of IPv6 addresses

- **Convention:** The 128-bit IPv6 address is written as **eight 16-bit integers** (using hexadecimal digits for each integer)

CEDF:BP76:3245:4464:FACE:2E50:3025:DF12

- **Short notation:**

- Abbreviations of leading zeroes:

CEDF:BP76:0000:0000:009E:0000:3025:DF12

→ CEDF:BP76:0:0:9E :0:3025:DF12

- “:0000:0000:0000” can be written as “::”

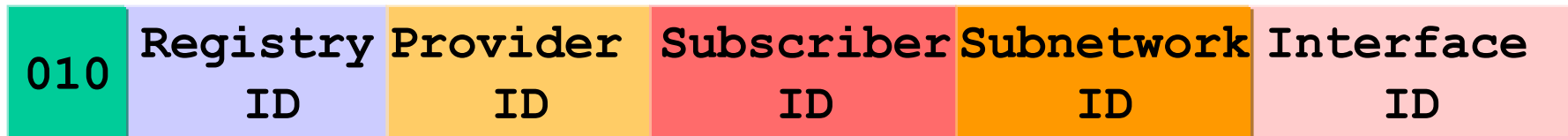
CEDF:BP76:0:0:FACE:0:3025:DF12 → CEDF:BP76::FACE:0:3025:DF12

- IPv6 addresses derived from IPv4 addresses have 96 leading zero bits. Convention allows to use IPv4 notation for the last 32 bits.

::80:8F:89:90 → ::128.143.137.144

IPv6 Provider-Based Addresses

- The first IPv6 addresses will be allocated to a provider-based plan



- **Type:** Set to “010” for provider-based addresses
- **Registry:** identifies the agency that registered the address
- The following fields have a variable length (recommended length in “()”)*
- **Provider:** Id of Internet access provider (16 bits)
- **Subscriber:** Id of the organization at provider (24 bits)
- **Subnetwork:** Id of subnet within organization (32 bits)
- **Interface:** identifies an interface at a node (48 bits)

More on IPv6 Addresses

- The provider-based addresses have a similar flavor as CIDR addresses
- IPv6 provides address formats for:
 - **Unicast** – identifies a single interface
 - **Multicast** – identifies a group. Datagrams sent to a multicast address are sent to all members of the group
 - **Anycast** – identifies a group. Datagrams sent to an anycast address are sent to one of the members in the group.